

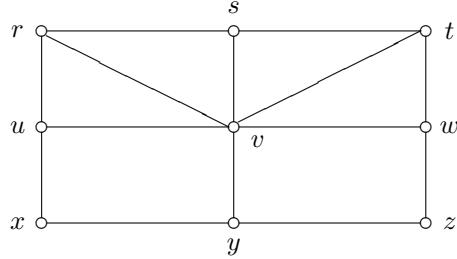
**Math 483 - Spring 26**

**Homework 1**

**SOLUTIONS**

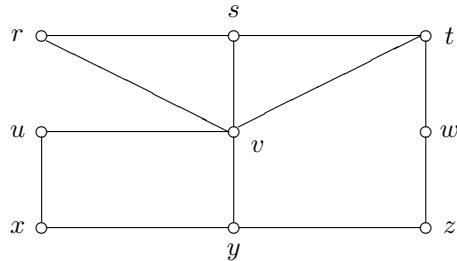
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1. Let  $G$  be the graph below:

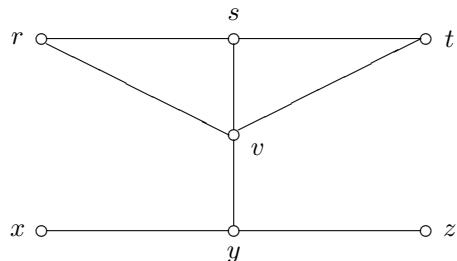


Let  $e$  and  $f$  be the edges  $e = ru$  and  $f = vw$ . Draw the subgraphs  $G - \{e, f\}$  and  $G - \{u, w\}$ .

**Answer.** Here is  $G = \{e, f\}$ :



Here is  $G - \{u, w\}$ :



2. For the graph above, give an example of each of the following, or explain why no such example exists:

(i) An  $x-y$  walk of length 6.

**Answer.** One possible example is  $W = (x, u, r, s, t, v, y)$ .

(ii) A  $v-w$  trail that is not a  $v-w$  path.

**Answer.** One possibility is  $T = (v, s, t, v, w)$ .

(iii) An  $r-z$  path of length 2.

**Answer.** There is no  $r-z$  path of length 2. Note that vertices  $w$  and  $y$  are the only vertices that are distance 1 from  $z$ ; and  $r$  is not adjacent to either  $w$  or  $y$ . So the distance from  $r$  to  $z$  has to be at least three, and there can be no path of length 2.

(iv) An  $x-z$  path of length 3.

**Answer.** There is no  $x-z$  path of length 3. To have an  $x-z$  path of length 3, we would need to have an path of length 2 from  $x$  to either  $y$  or to  $w$ ; but no such paths exist.

(v) An  $x$ - $t$  path of length  $d(x, t)$ .

**Answer.** Note that  $d(x, t) = 3$ . The vertices that are distance 1 from  $x$  are  $u$  and  $y$ . The vertices that are distance 2 are then  $r, v$ , and  $z$ . So the distance to  $t$  is at least 3. And since  $W = (x, y, v, t)$  is an  $x$ - $t$  path of length 3, we have that the distance is at most 3. Thus, the distance is 3, and  $W$  is an example of a path as desired.

(vi) A circuit of length 10.

**Answer.** One example could be  $C = (x, u, r, v, s, t, v, w, z, y, x)$ . Note that this is not a cycle, since the vertex  $v$  is repeated.

(vii) A cycle of length 8.

**Answer.** One possible example is  $C = (x, u, r, s, t, w, z, y, x)$ .

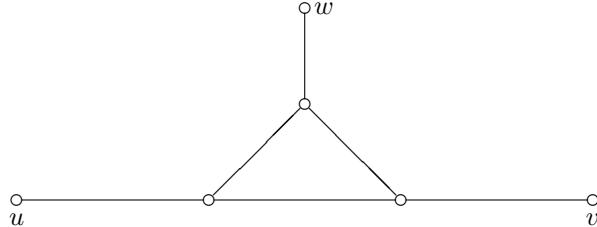
(viii) A geodesic of length  $\text{diam}(G)$ .

**Answer.** Note that the furthest apart that two vertices of  $G$  are from each other is 3: indeed, we can get from  $v$  to every vertex except  $x$  and  $z$  in one step; and from  $v$  to  $x$  and  $z$  in two steps. So we can get from  $x$  to any vertex other than  $z$  in at most three steps by going to  $v$  first; and we can get from  $x$  to  $z$  in two steps. Symmetrically, the distance from  $z$  to any vertex is at most three. And the distance between any two vertices other than  $x$  and  $z$  is at most two, again by going through  $v$  if necessary. And since we already know that the distance from  $x$  to  $t$  is 3, we have that the diameter of  $G$  is exactly 3.

So what we need is a geodesic of length 3. One possibility is  $P = (x, y, v, t)$ .

3. Give an example of a connected graph  $G$  that includes three pairwise distinct vertices  $u, v$ , and  $w$  (and possibly other vertices), such that  $d(u, v) = d(u, w) = d(v, w) = \text{diam}(G) = 3$ .

**Answer.** Here is an example:



4. We saw in class that a graph  $G$  of order at least 3 is connected if and only if  $G$  contains two distinct vertices  $u$  and  $v$  such that both  $G - u$  and  $G - v$  are connected. One might guess the following statement would also be true:

*Every connected graph  $G$  of order at least 4 contains three pairwise distinct vertices  $u, v$ , and  $w$ , such that each of  $G - u$ ,  $G - v$ , and  $G - w$  are connected.*

However, this statement is false. Draw a counterexample; that is, a graph  $G$  with at least 4 vertices, that is connected, and show that no choice of three distinct vertices will satisfy the condition.

**Answer.** One example of such a graph would be a path of length 3:



Any choice of three pairwise distinct vertices must include either  $v_2$  or  $v_3$ , and both  $G - v_2$  and  $G - v_3$  are disconnected. Yet  $G$  has order 4 and is connected.

5. Let  $G$  be a connected graph, and suppose that  $u$  and  $v$  are two distinct vertices in  $G$ . What is the minimum order of a connected subgraph of  $G$  that contains both  $u$  and  $v$ ? Explain your answer.

**Answer.** The minimum order is  $d(u, v) + 1$ , with a geodesic from  $u$  to  $v$  being the subgraph. Indeed,  $H$  is a connected subgraph that contains both  $u$  and  $v$ , then any  $u$ - $v$  path in  $H$  is also a  $u$ - $v$  path in  $G$ , and so contains at least  $d(u, v)$  edges, and hence  $d(u, v) + 1$  vertices. So the order of  $H$  is at least  $d(u, v) + 1$ . On the other hand, if  $W = (u = v_0, v_1, \dots, v_k = v)$  is a  $u$ - $v$  geodesic, so  $k = d(u, v)$ , then  $G[\{v_0, \dots, v_k\}]$  is a connected graph that contains exactly  $k + 1$  vertices, including  $u$  and  $v$ , so the minimum order is at most  $d(u, v) + 1$ . This proves the equality.